

**Library of Congress Cataloging-in-Publication Data**

Simon, Marvin Kenneth.  
Digital communication techniques / Marvin K. Simon, Sami M.  
Hinedi, William C. Lindsey.  
p. cm.  
Includes bibliographical references and index.  
Contents: v. 1. Signal design and detection  
ISBN 0-13-200610-3  
1. Digital communications. I. Hinedi, Sami M. II. Lindsey,  
William C. III. Title.  
TK5103.7.S52 1994 93-9053  
621.382—dc20 CIP

Editorial/production supervision and interior design: *Harriet Tellem*  
Cover design: *Aren Graphics*  
Manufacturing buyer: *Alexis Heydt*  
Acquisitions editor: *Paul W. Becker*  
Editorial assistant: *Maureen Diana*



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Prentice-Hall, Inc.  
A Paramount Communications Company  
Englewood Cliffs, New Jersey 07632

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Printed in the United States of America  
10 9 8 7 6 5 4 3 2 1

**ISBN: 0-13-200610-3**

Prentice-Hall International (UK) Limited, *London*  
Prentice-Hall of Australia Pty. Limited, *Sydney*  
Prentice-Hall Canada Inc., *Toronto*  
Prentice-Hall Hispanoamericana, S.A., *Mexico*  
Prentice-Hall of India Private Limited, *New Delhi*  
Prentice-Hall of Japan, Inc., *Tokyo*  
Simon & Schuster Asia Pte. Ltd., *Singapore*  
Editora Prentice-Hall do Brasil, Ltda., *Rio de Janeiro*

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### 2.2.3 Biphase (Manchester) Baseband Signaling

Here the two elementary signals are defined by

$$s_1(t) = \begin{cases} A; & 0 \leq t \leq T/2 \\ -A; & T/2 \leq t \leq T \end{cases} \quad (2.66)$$

$$s_2(t) = -s_1(t)$$

Substituting the Fourier transform of (2.66) into (2.59) gives

$$\begin{aligned} \frac{S_m(f)}{E} &= \frac{1}{T}(1-2p)^2 \sum_{\substack{n=-\infty \\ n \neq 0}}^{\infty} \left(\frac{2}{n\pi}\right)^2 \delta\left(f - \frac{n}{T}\right) \\ &+ 4p(1-p) \left[ \frac{\sin^4(\pi f T/2)}{(\pi f T/2)^2} \right] \end{aligned} \quad (2.67)$$

For  $p = 1/2$ , the line spectrum disappears and

$$\frac{S_m(f)}{E} = \frac{\sin^4(\pi f T/2)}{(\pi f T/2)^2} \quad (2.68)$$

### 2.2.4 Delay Modulation or Miller Coding

As indicated by Hecht and Guida [3], the Miller coding scheme can be modeled as a Markov source with four states whose stationary probabilities are all equal to  $1/4$  and whose transition matrix is given by

$$\mathbf{P} = \begin{pmatrix} 0 & 1/2 & 0 & 1/2 \\ 0 & 0 & 1/2 & 1/2 \\ 1/2 & 1/2 & 0 & 0 \\ 1/2 & 0 & 1/2 & 0 \end{pmatrix} \quad (2.69)$$

Another property of the Miller code is that it satisfies the recursion relation

$$\mathbf{P}^{4+i} \boldsymbol{\Gamma} = -\frac{1}{4} \mathbf{P}^i \boldsymbol{\Gamma} \quad i \geq 0 \quad (2.70)$$

where  $\boldsymbol{\Gamma}$  is the signal correlation matrix whose  $ik^{\text{th}}$  element is defined by

$$\gamma_{ik} \triangleq \frac{1}{\sqrt{E_i E_j}} \int_0^T s_i(t) s_k(t) dt \quad i, k = 1, 2, 3, 4 \quad (2.71)$$

For the Miller code, the four elementary signals are defined by

$$\begin{aligned} s_1(t) &= -s_4(t) = A; \quad 0 \leq t \leq T \\ s_2(t) &= -s_3(t) = \begin{cases} A; & 0 \leq t \leq T/2 \\ -A; & T/2 \leq t \leq T \end{cases} \end{aligned} \quad (2.72)$$

and  $E_i = A^2 T$ ;  $i = 1, 2, 3, 4$ . Substituting (2.72) into (2.71) and arranging the results in

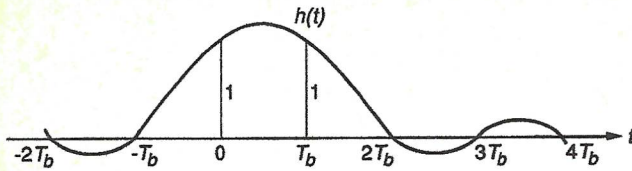


Figure 9.8 Duobinary pulse shape

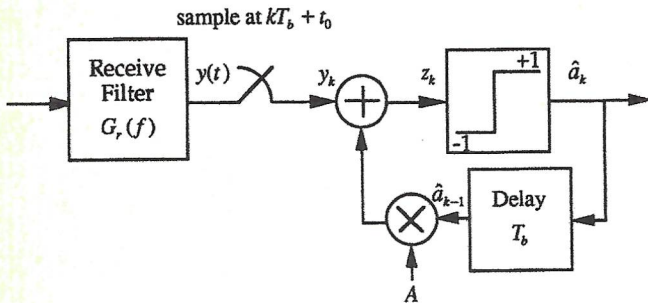


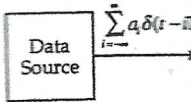
Figure 9.9 Duobinary receiver with perfect decision feedback

If we could always make *perfect* decisions, we could, in principle, subtract out the ISI and use a binary decision scheme with a single threshold. The appropriate receiver for doing this would be as shown in Fig. 9.9. If no error is made on the decision for  $a_{k-1}$ , that is,  $\hat{a}_{k-1} = a_{k-1}$ , then  $z_k = y_k - A\hat{a}_{k-1} = Aa_k + \tilde{n}_k$  and the probability of error for  $a_k$  would be  $\frac{1}{2} \operatorname{erfc} \left( A/\sqrt{2}\sigma_y \right)$  which is the identical result as that for binary Nyquist signaling. On the other hand, if an error was made on the decision for  $a_{k-1}$ , that is,  $\hat{a}_{k-1} = -a_{k-1}$ , then  $z_k = y_k - A\hat{a}_{k-1} = A(a_k + 2a_{k-1}) + \tilde{n}_k$ . Thus, depending on the relative polarities of  $a_k$  and  $a_{k-1}$ , the double ISI can help or hurt. The problem with this detection scheme then is that once an error occurs, it tends to *propagate* further errors, that is, if  $a_{k-1}$  is in error, then its effect is incorrectly compensated for when deciding  $a_k$ ; consequently,  $a_k$  is likely to be in error also.

To get around the problem of error propagation, Lender [6] suggested a means of avoiding the necessity for ISI compensation (subtracting out the single ISI) at the receiver. His scheme suggested that we *precode* the data  $\{a_i\}$  at the transmitter into the sequence  $\{d_i\}$  as follows. If we think of the binary data as 0's and 1's, then the precoder is described by  $d_n = a_n \oplus d_{n-1}$  where the symbol  $\oplus$  denotes modulo-2 addition. Note that  $a_n$  now indicates the *change* in the output symbol  $d_n$ . If  $a_n = 0$ , then there is no change in  $d_n$ . If  $a_n = 1$ , then there is a change in  $d_n$ . This simple encoding scheme is identical to differential encoding described in Chapter 7 in connection with differential detection of BPSK modulation. The precoder can also be reinterpreted in terms of "1"'s and "-1"'s for the input bits  $\{a_n\}$ . In particular, the modulo-2 summer is replaced by a multiplier resulting in the relation  $d_n = a_n \times d_{n-1}$ . Note that  $d_n$  is still a binary random sequence with equiprobable 1's and -1's.

Using the  $\pm 1$  representation for the  $a_n$ 's, we obtain the implementation for the duobinary system illustrated in Fig. 9.10. The mean of  $y_k$ , the sample of  $y(t)$  at  $t = kT_b + t_0$ , is now equal to  $A[d_k + d_{k-1}]$ . Table 9.2 gives the values for this mean in relation to the possible values of the input (information) binary symbol  $a_k$ .

Thus, we see that when the information bit is a "-1", then the detected signal sample is "0" (independent of  $a_{k-1}$ ). When the information bit is a "+1", then the detected signal



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